

Additional Exercises: Logics and Statistics for Language Modeling 2009-2010

1 First-Order Logic

- Construct a single first-order model that satisfies the following formulas:

$$\begin{aligned}\forall xyz.((R(x, y) \wedge R(y, z)) \rightarrow R(z, z)) \\ \exists xy.(P(x) \wedge \neg P(y)) \\ \forall x.\exists y.(R(x, y) \wedge x \neq y)\end{aligned}$$

- Construct a single first-order model that satisfies the following formulas:

$$\begin{aligned}\forall x.\exists y.R(x, y) \\ \exists xy.x \neq y \\ \forall xy.(R(x, y) \rightarrow \neg R(y, x)) \\ \forall x.(P(x) \vee Q(x)) \\ \forall x.(P(x) \rightarrow \neg \exists y.R(x, y))\end{aligned}$$

- Construct a single first-order model that satisfies the following formulas:

$$\begin{aligned}\exists xyz.(\text{Man}(x) \wedge \text{Woman}(y) \wedge \text{MarriedWith}(x, y) \wedge \text{DogOf}(x, z)) \\ \forall xy.(\text{MarriedWith}(x, y) \rightarrow \text{MarriedWith}(y, x)) \\ \forall xy.(\text{DogOf}(x, y) \rightarrow \text{Dog}(y)) \\ \forall x.(\text{Dog}(x) \rightarrow \text{has}(x, \text{tail}(x))) \\ \forall x.(\text{Dog}(x) \rightarrow \exists y.\text{Loyal}(x, y))\end{aligned}$$

- Construct a single first-order model that satisfies the following formulas:

$$\begin{aligned}\exists xyz.(\text{Man}(x) \wedge \text{Tall}(x) \wedge \text{MarriedWith}(x, y) \wedge \text{DogOf}(y, z)) \\ \forall xy.(\text{MarriedWith}(x, y) \rightarrow \text{MarriedWith}(y, x)) \\ \forall xyz.((\text{DogOf}(x, z) \wedge \text{MarriedWith}(x, y)) \rightarrow \text{DogOf}(y, z)) \\ \forall xy.(\text{DogOf}(x, y) \rightarrow \text{Dog}(y)) \\ \forall x.(\text{Dog}(x) \rightarrow \text{has}(x, \text{tail}(x)))\end{aligned}$$

- Try to unify the following terms, where we use capital letters for variables (all the others are functions or constants), and show the unifier substitution when successful.

1. $a(b, c) =^? a(X, Y)$
2. $a(X, c(d, X)) =^? a(2, c(d, Y))$
3. $a(X, Y) =^? a(b(c, Y), Z)$
4. $tree(left, root, Right) =^? tree(left, root, tree(a, b, tree(c, d, e)))$
5. $p(X, g(X), Y) =^? p(a, g(X), b)$
6. $p(1, 2) =^? q(X, Y)$
7. $p(a, X) =^? p(X, g(X))$
8. $p(X) =^? p(g(X))$
9. $p(h(X)) =^? p(g(X))$
10. $q(X, Y) =^? q(Y, X)$
11. $q(1, Y) =^? q(Y, X)$
12. $q(g(X, Y), f(X)) =^? q(g(h(Z), Y), f(h(Z)))$
13. $q(g(X, Y), f(X)) =^? q(g(h(a), Y), f(h(Z)))$
14. $p(f(X, g(a)), g(g(Y))) =^? p(f(X, Z), g(Z))$
15. $p(f(X), Z) =^? p(Y, z)$
16. $r(f(X, X)) =^? r(f(g(Y), Y))$
17. $app(cons(X, Y), nil) =^? app(Z, Y)$
18. $p(X, Z) =^? p(f(Z), a)$

- Show that φ does not entails ψ by producing, in each case, a model where the left formula is true and the right formula is false.

φ		ψ
$\exists x.P(x, x)$	\neq	$P(a, b)$
$\exists x.P(x, x)$	\neq	$P(a, a)$
$\forall x.P(f(x), g(x))$	\neq	$P(f(a), a)$
$\forall x.P(f(x), f(g(x)))$	\neq	$P(f(a), f(g(b)))$

- Given the following set of formulas:

$$\begin{aligned}
&\forall xy.(R(x, y) \vee R(y, x)) \\
&\exists xy.(P(x) \wedge \neg P(y)) \\
&\forall xy.(P(x) \rightarrow \neg R(x, y)) \\
&\forall xy.(R(x, y) \rightarrow R(y, x))
\end{aligned}$$

- Apply Skolemization and obtain the corresponding set of clauses.
- Use resolution to prove that the set is inconsistent.

- Given the following set of formulas:

$$\begin{aligned}
&\forall xy.(R(x, y) \rightarrow (P(y) \vee Q(y))) \\
&\forall xy.(R(x, y) \rightarrow \neg P(y)) \\
&\exists xy.(R(x, y) \wedge \neg Q(y))
\end{aligned}$$

- Apply Skolemization and obtain the corresponding set of clauses.
- Use resolution to prove that the set is inconsistent.